# OPTIMAL PRODUCTION OF BRA BLOCK FACTORY BY USING SIMPLEX METHOD 

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#### Abstract

In this paper simplex method is used to obtain the optimal solution to maximize the production profit for three different block types in BRA BLOCK FACTORY for heat insulating light construction pomic blocks productions. The three different model that we have taken to investigat the maximum profit for the factory are BB6H.10, BBS. 15 and BB2H.20. As aresult we found that the first and the third types of model gives the maximum profit. KEYWORDS : Linear Programming Problems;Integer Programming Problems; Simplex Method; Constraints; Objective Function; Variables, Slack Variable;Surplus Variables; Optimal Solution.


## 1. INTRODUCTION

Simplex method is the first method which was able to solve a linear programming problem developed by the American Air Force during the Second World War. It has proved in practice to perform very well in problems of small or medium size, although (Klee,V., Minty, GJ., 1972) had proven that in the worst case its behavior presents exponential complexity. Since then, a lot of research has been done to 1nd a faster (polynomial) algorithm that can solve LPs. The main branch of this research has been devoted to interior point methods (IPM). There is still much research being done in order to improve pivoting algorithms. Recently, (Paparrizos, 1991) proposed an exterior point Simplex algorithm that avoids the feasible Region; (Luh, H., Tsaih, R. , 2002) proposed using IPM to replace phase I of the Simplex; Paparrizos et al. In this research simplex method is used to find optimal solution for the problem that has taken from BRA.BLOCK FACTORY Duhok, three types of block have been chosen of

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PONZA model BB6H.10, BBS. 15 and BB2H.20, in this
paper $\mathbf{B}_{1}, \mathbf{B}_{2}$ and $\mathbf{B}_{3}$ represent the our chosen block types respectively. Furthermore, as it is clear that the linear programing problem is organized in objective function and constraints, in this paper our objective function is the percentage profit in each block types and the constraint part includes five constraints the first two one are mixing and pressing the others constraints are number of our block product in per day.
2. Linear programming problem (Lpp) (Usha, K., Rashmi , M., 2016), (Bland R. , 1977)
A linear programming problem (LPP) is a mathematical program in which the objective function is linear in the unknowns and the constraints consist of linear inequalities.
We formulate a mathematical model for general problem of allocating resources to activities. In particular, this model is to select the values for $\mathrm{x}_{1}, \mathrm{x}_{2}$, $\ldots . ., \mathrm{x}_{\mathrm{n}}$ as to maximize or minimize

$$
\mathrm{z}=\mathrm{c}_{1} \mathrm{x}_{1}+\mathrm{c}_{2} \mathrm{x}_{2}+\cdots+\mathrm{c}_{\mathrm{n}} \mathrm{x}_{\mathrm{n}}
$$

Subject to constraints :

$$
\begin{gathered}
a_{11} x_{1}+a_{12} x_{2}+\cdots+a_{1 n} x_{n}(\leq \text { or } \geq) b_{1} \\
a_{21} x_{1}+a_{22} x_{2}+\cdots+a_{2 n} x_{n}(\leq \text { or } \geq) b_{2} \\
\vdots \\
a_{n 1} x_{1}+a_{n 2} x_{2}+\cdots+a_{n n} x_{n}(\leq \text { or } \geq) b_{n}
\end{gathered}
$$

and

$$
\mathrm{x}_{1}, \mathrm{x}_{2}, \ldots . \mathrm{x}_{\mathrm{n}} \geq 0
$$

Where:
$\mathrm{Z}=$ value of overall measure of performance.
$x_{j}=$ level of activity ( for $j=1,2, \ldots \ldots, n$ )
$c_{j}=$ increase in Z that would result from unit increase in level of activity j .
$b_{i}=$ amount of resource ithat is available for allocation
to activities (( for $i=1,2, \ldots \ldots, m)$
$a_{i j}=$ amount of resource I consumed by each unit of activity j.

## 3. History of BRA Block factory

BRA plants was established on 02/02/2006 to produce the insulating volcanic light blocks, and the actual operation as on 12/01/2013 to be the largest plant of its kind in the Middle East, the production capacity of up to $(100,000$ pieces) a day $\ldots$ and nearly ( $2,000,000$ Piece) monthly using the technology of ADLER French company....
The plant produces more than ( 25 types) of BRA BLOCK with various dimensions, measurements according to German, European and Turkish standards
and specifications for this type of products ....
The production process are $100 \%$ automatic using volcanic material $100 \%$ natural without any additives or improvers, the mode of production and drying is done based on the principle of (natural curing system)

BRA BLOCK is produced from a natural volcanic stones which characterized by light weight density ranging between $600.700 \mathrm{~kg} / \mathrm{m} \mathrm{3}$, and high permeability to contain pores formed by natural interactions during the formation of these stones are called (PONZA) .....

4. Computational Procedure of Simplex Method (Dantzig, 1963)
Simplex method by G. Danztig in 1947. The simplex method provides an algorithm (a rule of procedure usually involving repetitive application of a prescribed operation) which is based on the fundamental theorem of linear programming.
This paper aims to find the maximum profit for three
types of ponza block which denoted by $B_{1}, B_{2}$ and $B_{3}$ of BRA. BLOCK FACTORY. Every type needs two stages, mixing and pressing,each step needs a period of time to produce a certain type of block. For this porpuse data are collected as shown in table1 and table 2.

### 4.1Types and amount of production for BRA Block

Table (1) : Available amont of production per day

| Block type | Amount available for each <br> day |
| :---: | :---: |
| Type $\left(\mathbf{B}_{1}\right)$ <br> $\left(\left(100^{*} 185 * 390 \mathrm{~mm}\right), 4 \mathrm{~kg}(\mathrm{dry}\right.$ weight $)$, <br> $\left.12.5 \mathrm{pcs} . / \mathrm{m}^{2}\right)$ | 33000 |
| Type $\left(\mathbf{B}_{2}\right)$ <br> $((150 * 185 * 4000 \mathrm{~mm}), 10 \mathrm{~kg}(\mathrm{dry}$ weight $)$, <br> $\left.12.5 \mathrm{pcs} . / \mathrm{m}^{2}\right)$ | 24000 |
| Type $\left(\mathbf{B}_{3}\right)$ <br> $\left(\left(200 * 185^{*} 4000 \mathrm{~mm}\right), 7 \mathrm{~kg}(\mathrm{dry}\right.$ weight $)$, <br> $12.5 \mathrm{pcs} . / \mathrm{m}^{2)}$ | 18000 |

### 4.2 Duration times of produce

In this stage the block is mixed and pressed in a different period of time as shown in table 2,
Table (2) : The Time needs in each stage

| Thages | $\mathbf{B}_{1}$ |  | $\mathbf{B}_{\mathbf{2}}$ |
| :---: | :---: | :---: | :---: |
| Types | $\mathbf{B}_{3}$ |  |  |
| mixing | 0.627 | 0.866 | 1.155 |
| pressing | 0.133 | 0.184 | 0.25 |
| Available time(per <br> hour) | 25200 | 25200 | 25200 |

Now the factory produces 33000 of type $B_{1}, 24000$ of type $B_{2}$ and 18000 of type $B_{3}$ per day. The profit of each product is $\$ 0.10, \$ .0 .05$ and $\$ 0.10$ respectively.

Table (3) : The profit of each types in US Dollar

| Block Type | Daily Profit |
| :---: | :---: |
| $\mathrm{B}_{1}$ | $\$ 3300$ |
| $\mathrm{~B}_{2}$ | $\$ 1200$ |
| $\mathrm{~B}_{3}$ | $\$ 1800$ |

5. Formulate the research problem as a linear minutes. The available times for both stages are not more programming problem
A BRA BLOCK FACTORY produces three types of blocks $B_{1}, B_{2}$ and $B_{3}$ and the profit of each product respectively are $\$ 0.10, \$ .0 .05$ and $\$ 0.10$. Each product is processes in two stages, mixing $S_{1}$ and pressing $S_{2}$, and needs 99 per second in $S_{1}$ stage and 21 per second in $S_{2}$ than 7 hours per day. The limit number of products $B_{1}, B_{2}$ and $B_{3}$ respectively are as follows : (330000, 240000, and 180000) per day. In this research we assume that $x_{1}$ represents the product number of type $B_{1}$, $x_{2}$ represents the product number of type $B_{2}$ and $x_{3}$ represents the product number of type $B_{3}$ stage. Available time for both stages is not more than 2

Table (4) : Formula

| Stages | Type B |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| $\mathbf{1}$ | Type B | 年 | Type $\mathbf{B}_{3}$ | Availability |
| Mixing $S_{1}$ | 0.627 | 0.866 | 1.155 | 25200 |
| Pressing $S_{2}$ | 0.133 | 0.184 | 0.250 | 25200 |
| Profit | $\$ 0.10$ | $\$ 0.05$ | $\$ 0.10$ |  |

The objective function for this problem is
Maximize $z_{x}=0.10 x_{1}+0.05 x_{2}+0.10 x_{3}$

The constraints are formulated as follows :

$$
0.627 x_{1}+0.866 x_{2}+1.155 x_{3} \leq 25200
$$

$0.133 x_{1}+0.184 x_{2}+0.250 x_{3} \leq 25200$

$$
x_{1} \leq 33000
$$

$$
\begin{equation*}
x_{2} \leq 24000 \tag{5.1}
\end{equation*}
$$

$$
x_{3} \leq 18000
$$

$$
\begin{array}{ll}
s_{1}=25200 & 0.627 x_{1}+0.866 x_{2}+1.155 x_{3}+ \\
s_{2}=25200 & 0.133 x_{1}+0.184 x_{2}+0.250 x_{3}+
\end{array}
$$

$$
x_{1}, x_{2}, x_{3} \geq 0
$$

First we start to convert the problem as standard linear programming problem, thus the slack variables should add in each constraint in terms of index (i.e. $s_{1}$ in constraint one, $s_{2}$ in constraint two and so on)
Objective function :
Maximize $z_{x}=0.10 x_{1}+0.05 x_{2}+0.10 x_{3}+0 s_{1}+0 s_{2}+$ $0 s_{3}+0 s_{4}+0 s_{5}$
Constraints
Table (5) : Table of initial data (Simplex Method)

| $\mathrm{C}_{\mathrm{j}}$ |  | 0.10 |  | 0.05 | $0.10 \quad 0$ | 0 | 0 |  | 00 |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Basic Variables | $\mathrm{C}_{\text {B }}$ | $\mathrm{X}_{\mathrm{B}}$ | $x_{1}$ | $x_{2}$ | $x_{3}$ | $s_{1}$ | $s_{2}$ | $s_{3}$ | $s_{4}$ | $s_{5}$ | Min ratio $\frac{\mathrm{X}_{\mathrm{B}}}{\mathrm{X}_{\mathrm{K}}}$ |
| $\mathrm{s}_{1}$ | 0 | 25200 | 0.627 | 0.866 | 1.155 | 1 | 0 | 0 | 0 | 0 | 40191.38755 |
| $\mathrm{s}_{2}$ | 0 | 25200 | 0.133 | 0.184 | 0.250 | 0 | 1 | 0 | 0 | 0 | 189473.68421 |
| $\mathrm{S}_{3}$ | 0 | 33000 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 33000 |
| $s_{4}$ | 0 | 24000 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | .... |
| $\mathrm{s}_{5}$ | 0 | 18000 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | .... |
| Z | $\mathrm{Z}=0$ |  | . 0.10 | . 0.05 | . 0.10 | 0 | 0 | 0 | 0 | 0 | $\longleftarrow \quad \Delta_{\mathrm{j}}$ |

In this step $\mathbf{s}_{\mathbf{3}}$ is removed and convert to $x_{1}$

| $\mathrm{C}_{\mathrm{j}}$ |  | $\longrightarrow$ | 0.10 |  | 0.05 | 0.10 |  | 0 | 0 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| BasicVariables | $\mathrm{c}_{\mathrm{B}}$ | $\mathrm{x}_{\mathrm{B}}$ | $x_{1}$ | $x_{2}$ | $x_{3}$ | $s_{1}$ | $s_{2}$ | $s_{3}$ | $S_{4}$ | $s_{5}$ | Min ratio $\frac{X_{B}}{X_{K}}$ |
| $\mathrm{s}_{1}$ | 0 | 4509 | 0 | 0.866 | 1.155 | 1 | 0 | .0.627 | 0 | 0 | 3903.8961 |
| $\mathrm{s}_{2}$ | 0 | 20811 | 0 | 0.184 | 0.250 | 0 | 1 | .0.133 | 0 | 0 | 83244 |
| $x_{1}$ | 0.10 | 33000 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | .... |
| $s_{4}$ | 0 | 24000 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | .... |
| $\mathrm{S}_{5}$ | 0 | 18000 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 18000 |
| Z | $\mathrm{Z}=3300$ |  | 0 | . 0.05 | . 0.10 | 0 | 0 | 0.10 | 0 | 0 | $\longleftarrow \quad \Delta_{\mathrm{j}}$ |

In this step $s_{1}$ is removed and convert to $x_{3}$

| $\mathrm{C}_{\mathrm{j}} \longrightarrow 0.10$ |  | 0.05 | 0.10 | 0 | 0 | 0 | 0 | 0 |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Basic <br> Variable <br> s | $\mathrm{C}_{\mathrm{B}}$ | $\mathrm{X}_{\mathrm{B}}$ | $x_{1}$ | $x_{2}$ | $x_{3}$ | $s_{1}$ | $s_{2}$ | $s_{3}$ | $s_{4}$ | $s_{5}$ | Min <br> ratio <br> $\frac{\mathrm{X}_{\mathrm{B}}}{\mathrm{x}_{\mathrm{K}}}$ |


| $x_{3}$ | 0.10 | 3903.8961 | 0 | 0.749783 | 1 | 0.8658 | 0 | .0 .5428571 | 0 | 0 |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathrm{~s}_{2}$ | 0 | 19835.0259 | 0 | .0 .0034457 | 0 | .0 .21645 | 1 | 0.002714275 | 0 | 0 |  |
| $x_{1}$ | 0.10 | 33000 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |  |
| $s_{4}$ | 0 | 24000 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 |  |
| $\mathrm{~s}_{5}$ | 0 | 14096.1039 | 0 | .0 .749783 | 0 | .0 .8658 | 0 | 0.5428571 | 0 | 1 |  |
| Z | $\mathrm{Z}=3690.38961$ | 0 | 0.0249783 | 0 | 0.08658 | 0 | 0.04571429 | 0 | 0 | $\leftarrow \Delta_{\mathrm{j}}$ |  |

Finally in above table we get the optimal solution for our problem which is $\mathrm{Z}=3690.38961$ where $x_{1}=$ $33000, x_{2}=0$ and $x_{3}=3903.8961$, here producing 33000 of type B1 and 3903.8961 of type B3 we get $\$ 3690.38961 \mathrm{amout}$ of money as maximum profit. In the above problem we take three types and find the maximum profit, but in case we need to find the maximum profit for two types, we will have three cases: Case 1
Hence for B1 and B3 we will get a result as the same as above result.

## Case 2

If the products of type B2 and product of type B3
Maximize $z_{x}=0.05 x_{2}+0.10 x_{3}$
Subject to :

$$
\begin{gathered}
0.866 x_{2}+1.155 x_{3} \leq 25200 \\
0.184 x_{2}+0.250 x_{3} \leq 25200 \\
x_{2} \leq 24000 \\
x_{3} \leq 18000
\end{gathered}
$$

$$
x_{2}, x_{3} \geq 0
$$

Using the same technique as used as in solving problem (5.1) the optimal solution $\mathrm{Z}=2054.61893764434$ where $x_{2}=5092.37875288683608$ and $x_{3}=18000 \mathrm{can}$ be obtained.

## Case 3

If the products of type B1 and products of type B2
Maximize $z_{x}=0.10 x_{1}+0.05 x_{2}$
Subject to :

$$
\begin{gathered}
0.627 x_{1}+0.866 x_{2} \leq 25200 \\
0.133 x_{1}+0.184 x_{2} \leq 25200 \\
x_{1} \leq 33000 \\
x_{2} \leq 24000 \\
x_{1}, x_{2} \geq 0
\end{gathered}
$$

Using the same technique as used as in solving problem (5.1) the optimal solution $Z=3560.33487297921$ where $x_{1}=33000$ and $x_{2}=5206.69745958429121 \mathrm{can}$ be obtained.

Table (6) : The maximum profit for defferent cases using simplix method

| Products | Production in seven hours per day |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Type B1 | 33000 | 0 | 0 | 33000 | 0 | 33000 |  |
| Type B2 | 0 | 24000 | 0 | 5206.69745958429121 | 5092.37875288683 | 0 |  |
| Type B3 | 0 | 0 | 18000 | 0 | 18000 | 3903.8961 |  |
| Profit \$ | 3300 | 1200 | 1800 | 3560.33487297921 | 2054.61893764434 | 3690.38961 |  |

Here we need to convert the results to be integer value in staid of rational number because in the real life the above results are not accepted, then the next step we need to use integer programming problem to meet our final aim.
6. Integer programming problem (IPP) (Schrijver, 1986) The linear programming models that have been discussed thus far all have been continuous, in the sense that decision variables are allowed to be fractional. Often this is a realistic assumption. For instance, we might easily produce $102 \frac{3}{4}$ gallons of a divisible good such as
wine. It also might be reasonable to accept a solution giving an hourly production of automobiles at $58 \frac{1}{2}$ if the model were based upon average hourly production, and the production had the interpretation of production rates.
At other times, however, fractional solutions are not realistic, and we must consider the optimization problem :
Maximize of objective function :

$$
z=\sum_{j=1}^{n} c_{j} x_{j}
$$

Subject to constraints :

$$
\begin{gathered}
\sum_{j=1}^{n} a_{i j} x_{j}=b_{i} \quad(i=1,2, \ldots, m) \\
x_{j} \geq 0 \quad(j=1,2, \ldots, n)
\end{gathered}
$$

$x_{j}$ integer for some or all $(j=1,2, \ldots, n)$.
This problem is called the linear integer programming problem.
It is said to be a mixed integer program when some, but
not all, variables are restricted to be integer, and is called a pure integer program when all decision variables must be integers. As we saw in the preceding paper, if the constraints are of a network nature, then an integer solution can be obtained by ignoring the integrality restrictions and solving the resulting linear program. In general, though, variables will be fractional in the linear.programming solution, and further measures must be taken to determine the integer.programming solution. First, solve the problem as linear programming problem to obtain optimal solution :
The technique that used in this paper is called Technique rounding solution for solving it. The purpose of using the mentioned technique is obtaining integer number of blocks as it is known that after the calculations sometimes, we get a rational number of objects that should be integer. Therefor in this case we have to change original number to integer because otherwise in real life it does have no meaning.

Table (7) : Technique rounding solution

| Type B1 | 33000 | 0 | 0 | 33000 | 0 | 33000 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Type B2 | 0 | 24000 | 0 | 5206 | 5092 | 0 |
| Type B3 | 0 | 0 | 18000 | 0 | 18000 | 3903 |
| Profit \$ | 3300 | 1200 | 1800 | 3560.30 | 2054.60 | 3690.30 |

## Conclusion

In conclusion, after defining A BRA BLOCK FACTORY and collecting data concerning of three types of products BB6H.10, BBS.15and BB2H.20, we formulate a problem of three variables representing the block types then used Simplex method and integer programming to find the maximum profit for factory of seven hours per day. As a result, we obtain that producing 33000 of type BB6H. 10 and 3903 of type BB2H.20, give the maximum profit. It means that when the factory produces the mentioned number of types BB6H. 10 and BB2H. 20 it will get $\$ 3690.30$ as a profit per day.

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